### Graph Algorithms

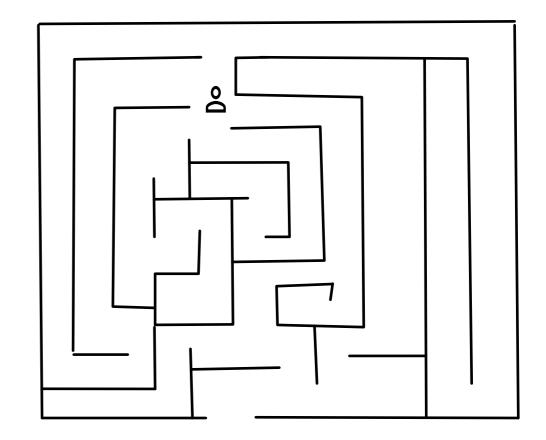
Thomas Schwarz, SJ

### Searching in Graphs

- Exploring a maze
  - You are in the middle of a maze
    - How do you get out
  - Ariadne's solution:
    - Use a thread of glittering jewels in order to avoid using the same edges several times
  - Follow a wall
    - Works for simple mazes

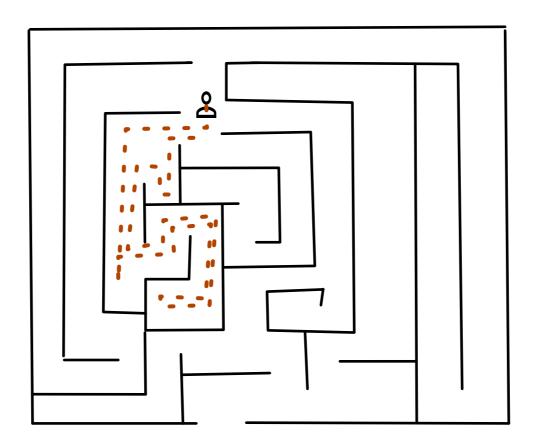


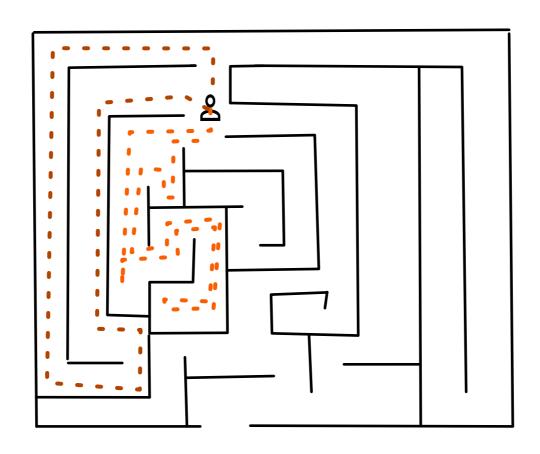
- Trémaux's Algorithm aka Hansel and Gretel's aka Ariadne's
  - Carry bread and leave bread crumbs on each path you follow
  - If you come to an intersection, follow one where there are no bread crumbs, if you can
  - If you come to an intersection and everything has already been marked or you are at a deadend, turn around if you came at a path that has only one thread of crumbs

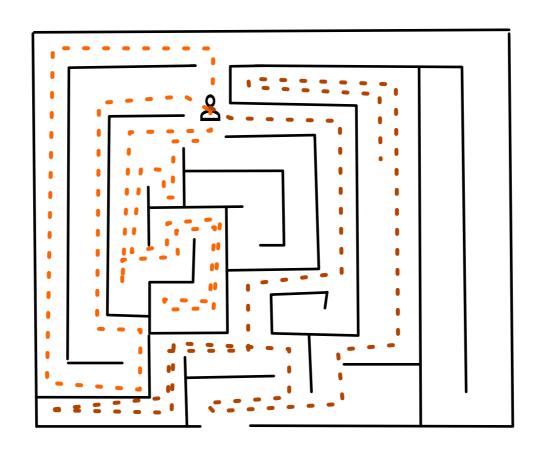


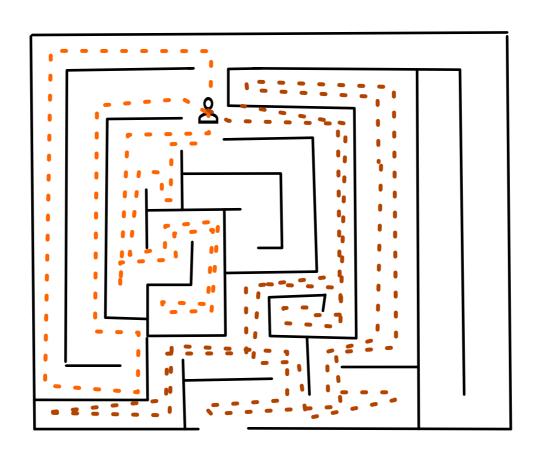
 If not, follow a path that has only one trail of crumbs.

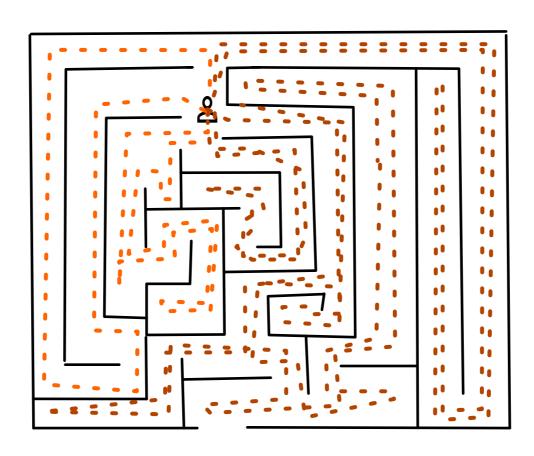
Example











https://en.wikipedia.org/wiki/
 File:Tremaux\_Maze\_Solving\_Algorithm.gif

- At the end:
  - All paths will be double marked and you will end up at the starting point
  - This means that you walked by the entry

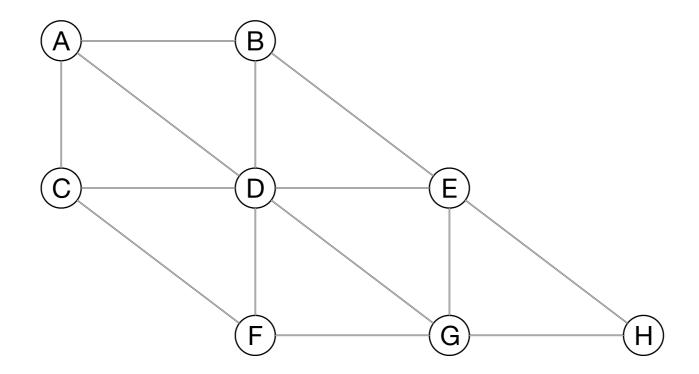
### Searching in Graphs

- We can use this idea for defining the first graph exploration algorithm.
  - Goal is to visit all vertices
  - We use a timer:
    - Starts out at 0
    - Incremented every time we do something
  - All nodes get marked with a
    - Discovery time: First time that we see the node
    - Finishing time: When we are done with the node

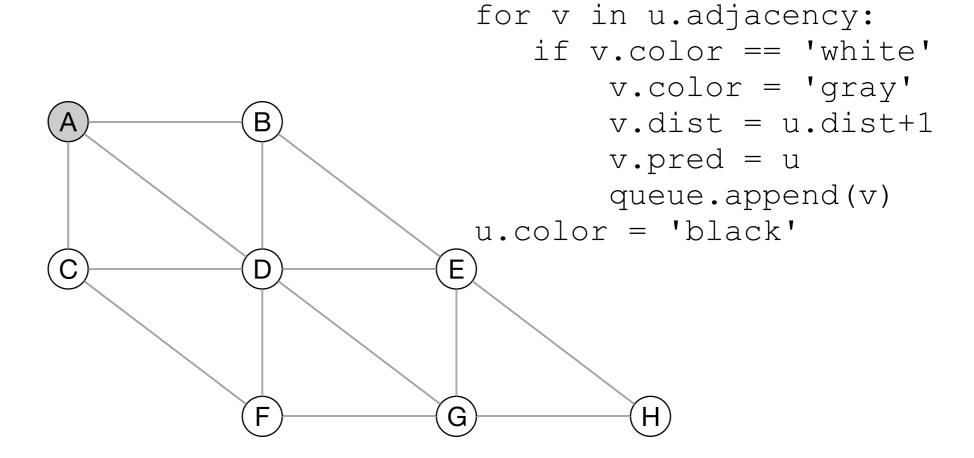
- Color a vertex
  - white: vertex has not yet been discovered
  - gray: vertex has been discovered, but still needs to be a base for exploration
  - black: vertex has been dealt with

```
bfs (G, s):
   for v in G. vertices:
       v.color = 'white'
       v.dist = inf
      v.pred = None
    s.color = 'gray'
    s.dist = 0
    s.pred = None
   queue = []
   queue.append(s)
   while queue:
       u = queue.pop(0)
       for v in u.adjacency:
          if v.color == 'white'
              v.color = 'gray'
              v.dist = u.dist+1
              v.pred = u
              queue.append(v)
       u.color = 'black'
```

• Example: s=A



- queue = { }
- u = A



while queue:

u = queue.pop(0)

• queue = {B,C,D}

```
while queue:
    u = queue.pop(0)
    for v in u.adjacency:
    if v.color == 'white'
        v.color = 'gray'
        v.dist = u.dist+1
        v.pred = u
        queue.append(v)
    u.color = 'black'
```

• queue = {B,C,D}

```
• queue = {C,D}

• u = B

• u = gueue.pop(0)

for v in u.adjacency:
    if v.color == 'white'
        v.color = 'gray'
        v.pred = u
        queue.append(v)
        u.color = 'black'
```

• queue = {C,D}

```
• queue = {C,D}

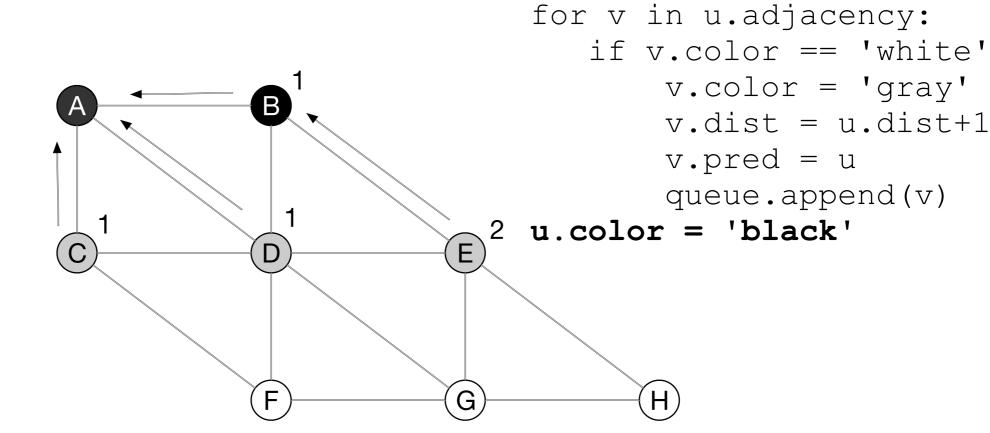
u = queue.pop(0)

for v in u.adjacency:
    if v.color == 'white'
    v.color = 'gray'
    v.dist = u.dist+1
    v.pred = u
    queue.append(v)

u.color = 'black'
```

• queue = {C, D, E}

- queue =  $\{C,D,E\}$
- u = B

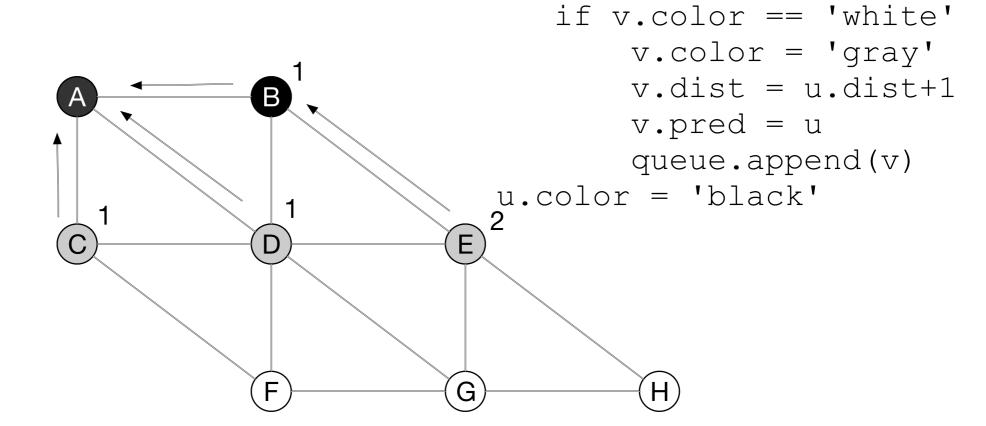


while queue:

u = queue.pop(0)

• queue = {C,D,E}

- queue =  $\{C,D,E\}$
- u = C



while queue:

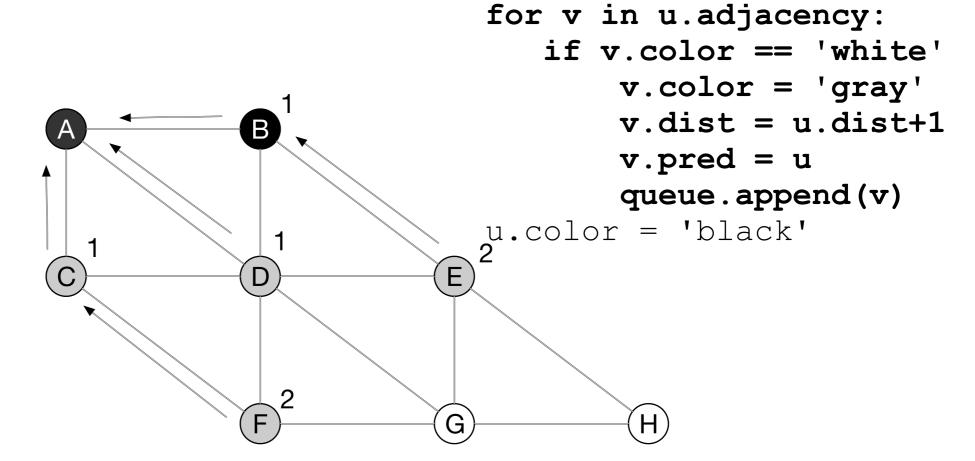
u = queue.pop(0)

for v in u.adjacency:

• queue = {D,E}

```
    queue = {D,E}
```

• u = C



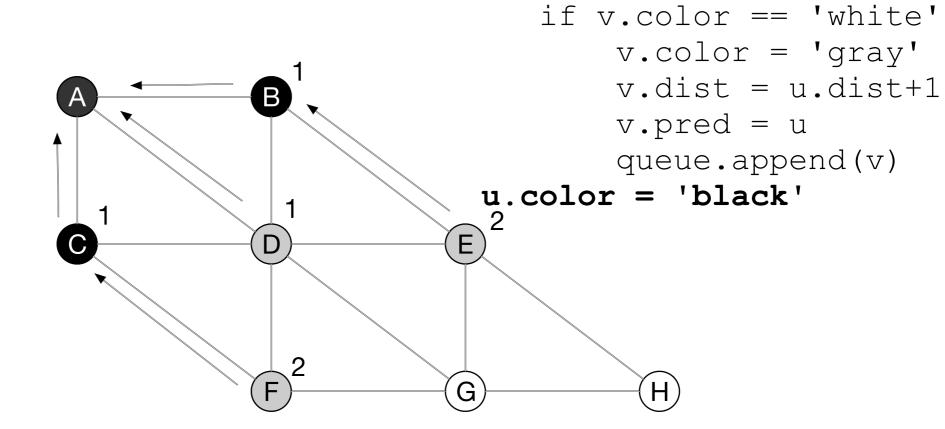
while queue:

u = queue.pop(0)

• queue = {D, E, F}

```
• queue = \{D,E,F\}
```

• u = C



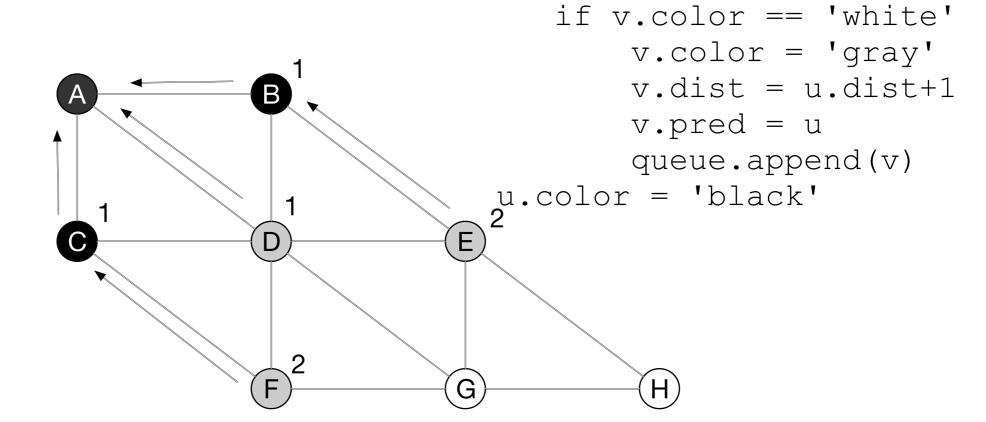
while queue:

u = queue.pop(0)

for v in u.adjacency:

• queue = {D, E, F}

- queue =  $\{D,E,F\}$
- u = D



while queue:

u = queue.pop(0)

for v in u.adjacency:

• queue = {E, F}

```
• queue = {E,F}

• u = D

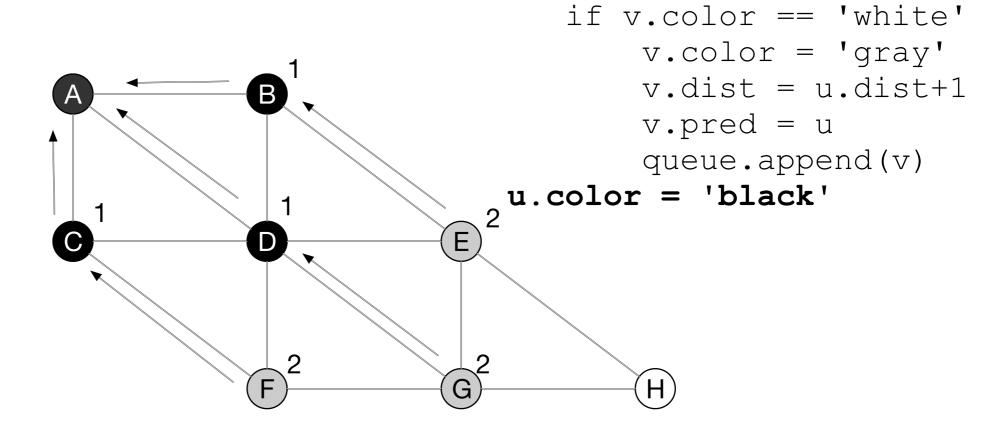
while queue:
    u = queue.pop(0)
    for v in u.adjacency:
    if v.color == 'white'
    v.color = 'gray'
    v.dist = u.dist+1
    v.pred = u
    queue.append(v)

2 u.color = 'black'
```

• queue = {E, F, G}

```
• queue = \{E,F,G\}
```

• u = D



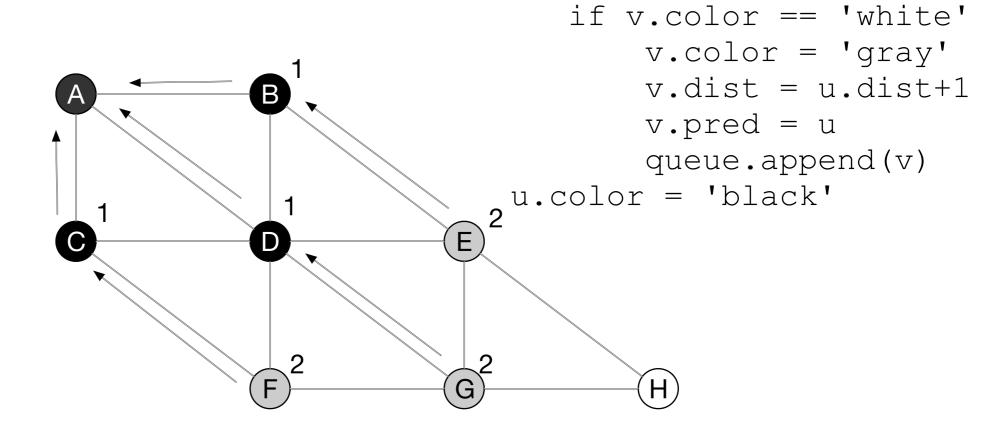
while queue:

u = queue.pop(0)

for v in u.adjacency:

• queue = {E, F, G}

- queue =  $\{E,F,G\}$
- u = E



while queue:

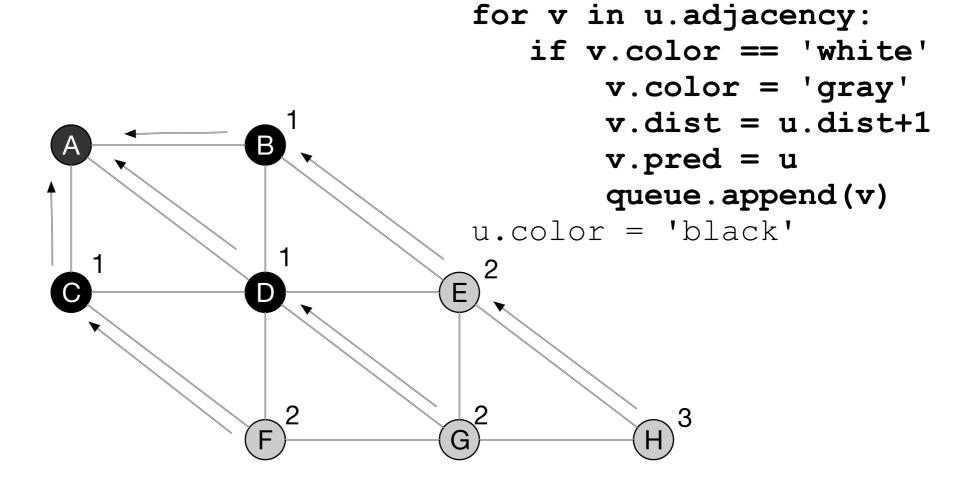
u = queue.pop(0)

for v in u.adjacency:

• queue = { F, G}

```
• queue = \{F,G\}
```

• u = E



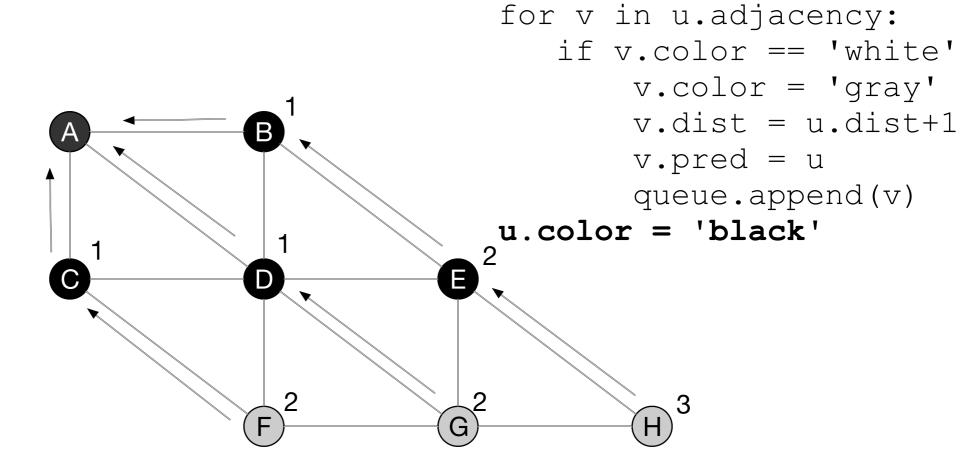
while queue:

u = queue.pop(0)

• queue = { F, G, H}

```
• queue = \{F,G,H\}
```

• u = E



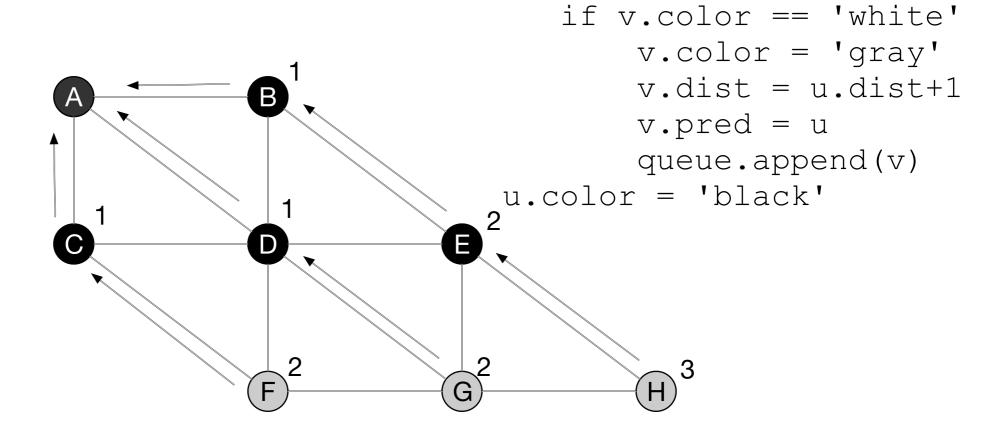
while queue:

u = queue.pop(0)

• queue = { F, G, H}

```
• queue = \{G,H\}
```

• u = F



while queue:

u = queue.pop(0)

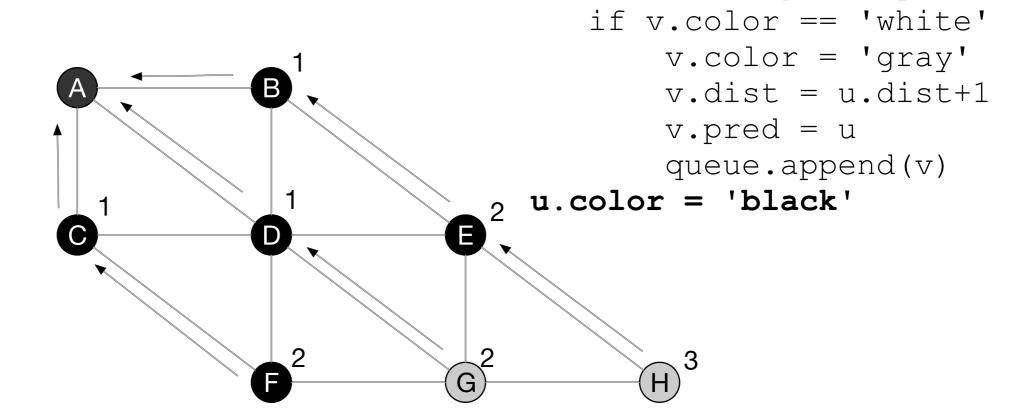
for v in u.adjacency:

• queue = {G, H}

• queue = {G, H}

```
• queue = \{G,H\}
```

• u = F



while queue:

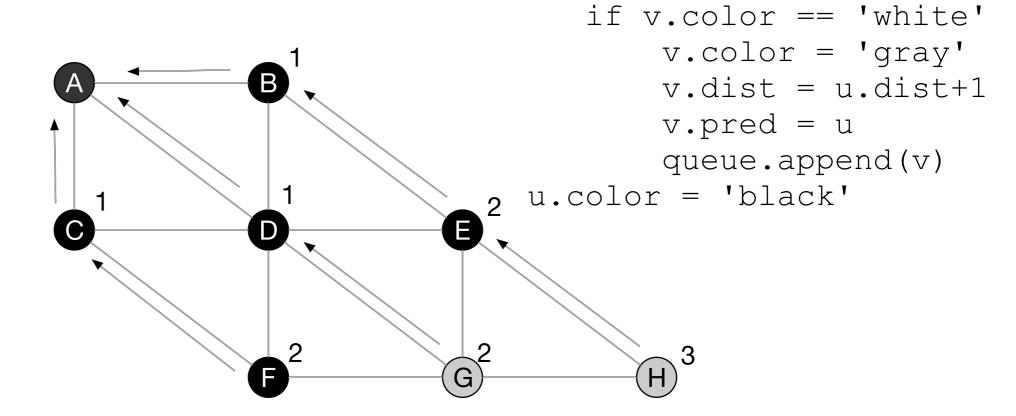
u = queue.pop(0)

for v in u.adjacency:

• queue = {G, H}

```
• queue = \{G,H\}
```

• u = G



while queue:

u = queue.pop(0)

for v in u.adjacency:

• queue = {H}

```
• queue = {G,H}

• u = queue.pop(0)
for v in u.adjacency:
    if v.color == 'white'
        v.color = 'gray'
        v.pred = u
        queue.append(v)
        u.color = 'black'
```

• queue = {H}

```
    queue = {H}
    u = queue for v in if v.
    u = G
```

```
u = queue.pop(0)
for v in u.adjacency:
    if v.color == 'white'
        v.color = 'gray'
        v.dist = u.dist+1
        v.pred = u
        queue.append(v)

u.color = 'black'
```

• queue = {H}

```
• queue = {}

• queue = {}

• u = queue.pop(0)

for v in u.adjacency:
    if v.color == 'white'
        v.dist = u.dist+1
        v.pred = u
        queue.append(v)
        u.color = 'black'
```

• queue = { }

```
while queue:
queue = { }
                                     u = queue.pop(0)
                                     for v in u.adjacency:
• u = H
                                        if v.color == 'white'
                                            v.color = 'gray'
                      B
                                            v.dist = u.dist+1
                                            v.pred = u
                                            queue.append(v)
                                   2 u.color = 'black'
```

• queue = { }

```
while queue:
queue = { }
                                     u = queue.pop(0)
                                     for v in u.adjacency:
• u = H
                                        if v.color == 'white'
                                            v.color = 'gray'
                      B
                                            v.dist = u.dist+1
                                            v.pred = u
                                             queue.append(v)
                                   2 u.color = 'black'
```

• queue = { }

- As you can see, BFS is just a version of Dijkstra's algorithm
- Distance calculates accurately the distance from the starting point
- The pred property allows us to generate a shortest path from the initial node
- We now prove these properties exactly

- Lemma: Let G = (E, V) be an undirected or directed graph. Let  $s \in V$  be an arbitrary vertex. Then for any edge  $(u, v) \in E$ 
  - $\delta(s, v) \le \delta(s, u) + 1$





• Recall:  $\delta(a,b)$  is the length of a shortest path from a to b

- Proof:
  - Assume first that  $\delta(s, u) = \infty$ , i.e. there is no path from s to u
    - Then  $\delta(s, v) \leq \infty = \delta(s, u) + 1$  regardless whether there is a path from s to v.

#### Proof:

- Next assume that  $\delta(s, u) < \infty$ , i.e. that there is a path from s to u.
  - Extend this path to a path from s to v.
  - This path has length  $\delta(s, u) + 1$ .
  - Then  $\delta(s, v) = \min(\text{Lenght of a path from } s \text{ to } v)$ 
    - ≤ Length of this path

• 
$$=\delta(s,u)+1$$

- Lemma: Let Let G=(E,V) be an undirected or directed graph. Let  $s\in V$  be an arbitrary vertex. Run BFS on G and s. Then for every vertex  $v\in V$ , v dist  $\geq \delta(s,v)$ .
  - This means that the calculated distance in BFS is at least as large as the actual distance

- Proof by induction on the number of enqueue operations
  - Notice that v.dist is assigned just when we are about to enqueue it

```
while queue:
    u = queue.pop(0)
    for v in u.adjacency:
        if v.color == 'white'
            v.color = 'gray'
- - >            v.dist = u.dist+1
            v.pred = u
            queue.append(v)
            u.color = 'black'
```

- Induction Start:
  - When s is enqueued all distance properties are infinity
    - with the exception of s which has dist 0
  - At this point, for every vertex  $v \in V$ , v . dist  $\geq \delta(s, v)$

- Induction step:
  - The value of the distance property only changes when we make the assignment just before enqueuing a white vector
  - Induction hypothesis implies u . dist  $\geq \delta(s, u)$
  - Therefore

```
v \cdot dist = u \cdot dist + 1 \ge \delta(s, u) + 1 \ge \delta(s, v)
```

```
while queue:
    u = queue.pop(0)
    for v in u.adjacency:
        if v.color == 'white'
            v.color = 'gray'
- - >       v.dist = u.dist+1
            v.pred = u
            queue.append(v)
    u.color = 'black'
```

 Afterwards, the vertex v is no longer white and never changes its distance value

- We now need to see more closely how the algorithm works:
  - We can think of the queue as the boundary between black and white vertices that moves slowly away from s
- Lemma: If the queue has vertices  $(v_1, v_2, ..., v_n)$  with  $v_1$  being the head, then
  - $v_n$  . dist  $\leq v_1$  . dist+1
    - and
  - $v_i$  . dist  $\leq v_{i+1}$  . dist for  $i=1,2,\ldots,n-1$

- Proof by induction on the number of queue operations
  - Initially, the queue has only s in it, so the property certainly holds
  - The queue changes through enqueuing and dequeuing operations

- If the head  $v_1$  is dequeued,  $v_2$  becomes the new head.
  - (If there is no  $v_2$  then the queue is empty, and the assertion holds vacuously)
  - Before dequeuing,  $v_1$  . dist  $\leq v_2$  . dist, therefore  $v_n$  . dist  $\leq v_1$  . dist  $+1 \leq v_2$  . dist
  - Therefore, the first inequality is true
  - The second assertion just looses the first inequality

• If a new element  $v_{n+1}$  is enqueued, we just dequeued a vertex u and are adding all white vertices adjacent to u

```
while queue:
    u = queue.pop(0)
    for v in u.adjacency:
        if v.color == 'white'
            v.color = 'gray'
- - >            v.dist = u.dist+1
            v.pred = u
            queue.append(v)
    u.color = 'black'
```

- Therefore,  $v_{n+1}$  . dist = u . dist + 1.
- By induction hypothesis, u . dist  $\leq v_1$  . dist because u and  $v_1$  were just in the same queue

- Therefore
  - $v_{r+1}$  . dist = u . dist+1  $\leq v_1$  . dist+1
- Proving the first assertion

- From the induction hypothesis, we also have
  - $v_n$  . dist  $\leq u$  . dist +1
- which implies that
  - $v_n$  . dist  $\leq u$  . dist  $+1 \leq v_{n+1}$  . dist
- This is the only new part of the second assertion

- Breadth first search uses a queue
  - In Python, a
     queue is a list to
     which you
     append and from
     which you pop
  - C++ and Java have libraries that implement queues

```
def bfs(G, s):
    for u in G. Vertices:
        u.color = "white"
        u.d = infty
        u.pred = Null
    s.color="gray"
    s.d = 0
    s.pred = Null
    queue = Queue.queue()
    queue.enqueue(s)
    while queue:
        u = queue.head()
        for v in u.adjacency list:
            if v.color=="white"
                v.color = "gray"
                v.d = u.d + 1
                v.pred = u
                queue.enqueue(v)
        u.color="black"
```

- Depth first search replaces the queue with a stack
  - This changes the behavior of the algorithm considerably
  - Remarkably, the resulting Depth First Search is the more important and interesting algorithm

- Depth first search
- Version 1
  - Change queue into stack
  - Get rid of the distance

```
def dfs(G, s):
    for u in G. Vertices:
        u.color = "white"
        u.d = infty
        u.pred = Null
    s.color="gray"
    s.pred = Null
    queue = Stack.stack()
    stack.push(s)
    while stack:
        u = stack.pop()
        for v in u.adjacency list:
            if v.color=="white"
                v.color = "gray"
                v.pred = u
                stack.push(v)
        u.color="black"
```

- We add visiting times to our nodes:
  - Discovered time
    - When a node turns gray
  - Finished time
    - When a node turns black
- Because
  - some derived algorithms use it
  - in order to argue about DFS
- Whenever we change a node color, we increment a clock

- Unlike BFS, a typical DFS will want to classify all nodes
  - Have a DFS\_Visit(start\_node) that starts in a node and visit what can be available
  - Have a DFS() function that uses the visit function repeatedly if necessary

```
dfs visit(u):
    global clock
    clock += 1
    u.d = clock
    u.color = 'gray'
    for each v in u.adjacency:
        if v.color == 'white'
            v.pred = u
            dfs visit(v)
    u.color = 'black'
    clock += 1
    u.f = clock
```

```
dfs(G):
    for vertex in G.V:
        vertex.color = 'white'
        vertex.pred = None
    global clock = 0
    for vertex in G.V:
        if vertex.color = 'white':
            dfs visit(vertex)
```

- Understanding the algorithm
  - The stack is hidden in the recursive call
  - We can unroll it
  - But need to be careful as something on the stack can be already found and processed via another route

```
Start with C
```

```
OS stack is dfs_visit(C)
```

```
dfs_visit(u):

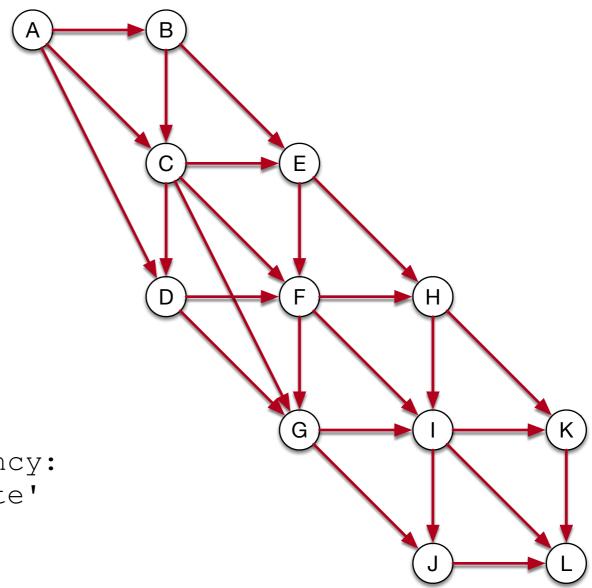
u.color = 'gray'

for each v in u.adjacency:

if v.color == 'white'

dfs_visit(v)

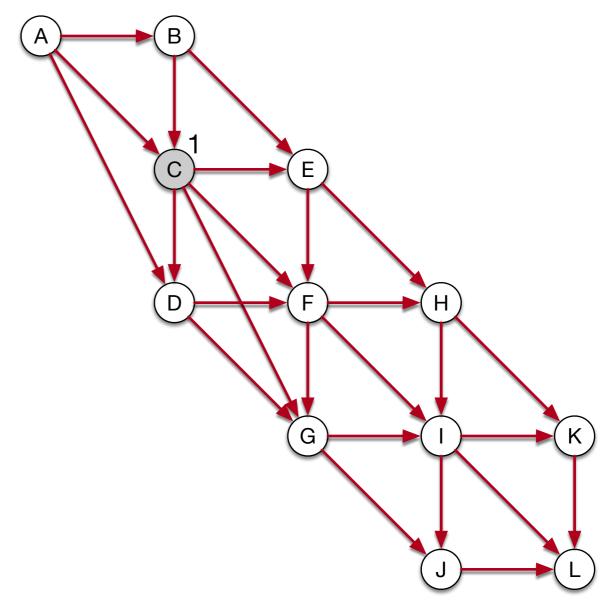
u.color = 'black
```



```
Start with C

OS stack

dfs_visit(C)
```



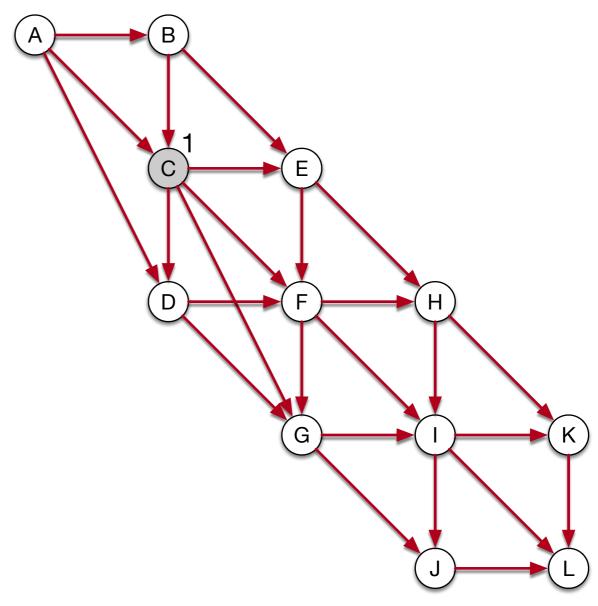
We set the clock to 1
We pick arbitrarily E
from the adjacency list

```
OS stack

dfs_visit(E)

dfs_visit(C)
```

```
dfs_visit(u):
    u.color = 'gray'
    for each v in u.adjacency:
        if v.color == 'white'
        dfs_visit(v)
        u.color = 'black
```

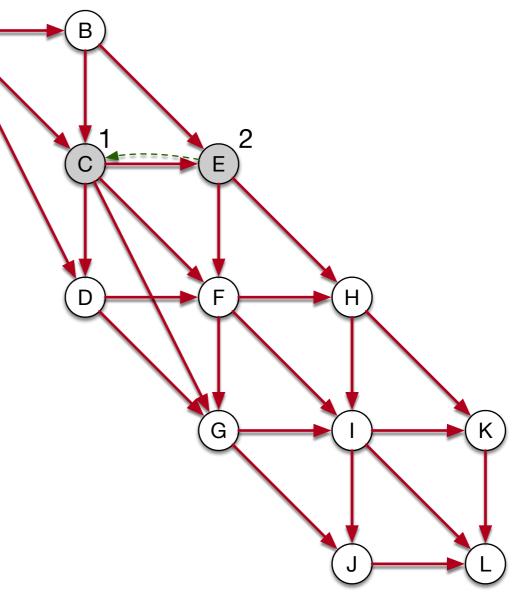


We call dfs\_visit(E)

```
OS stack

dfs_visit(E)

dfs_visit(C)
```

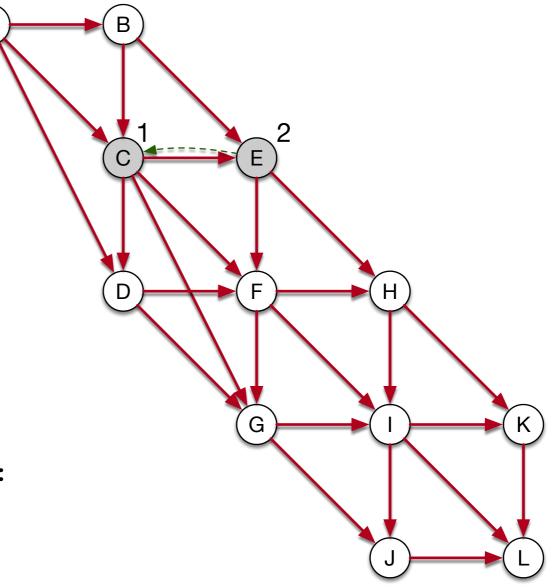


u = E, set E's discovery
time and set the
predecessor link

```
OS stack

dfs_visit(E)

dfs_visit(C)
```

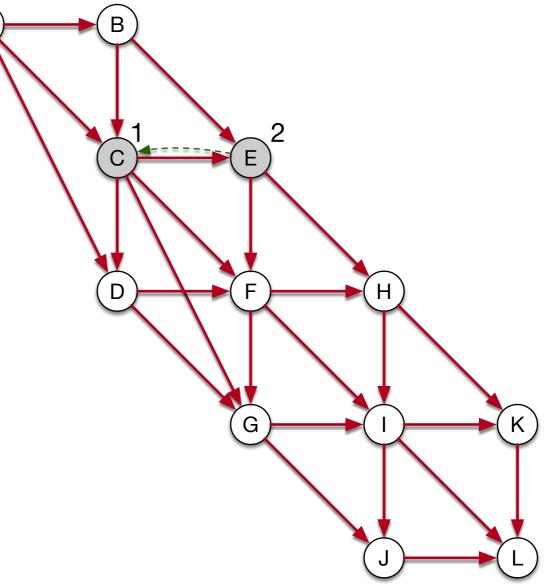


we pick v = H

```
OS stack

dfs_visit(E)

dfs_visit(C)
```



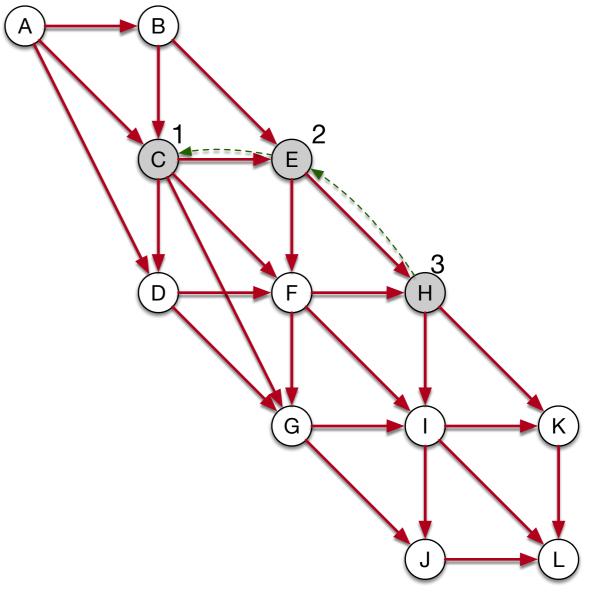
we pick v = H and call
dfs\_visit(H)

```
OS stack

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



we pick v = H and call
dfs\_visit(H)
this colors H gray

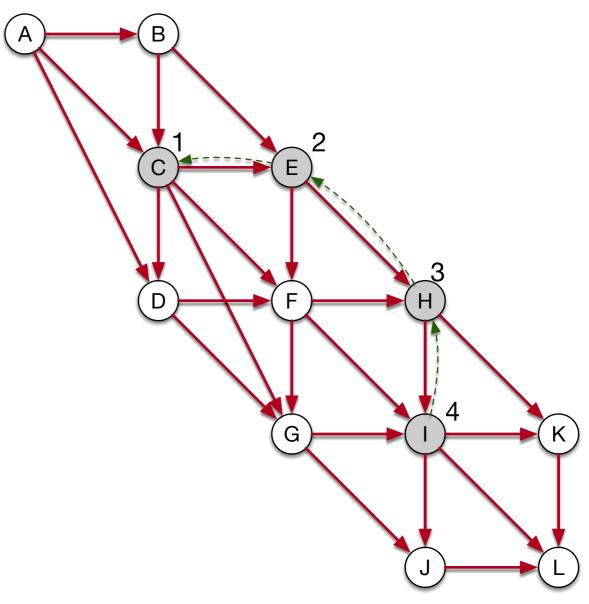
```
OS stack

dfs_visit(I)

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



we pick v = I and call
dfs\_visit(I)
this colors I gray

```
OS stack

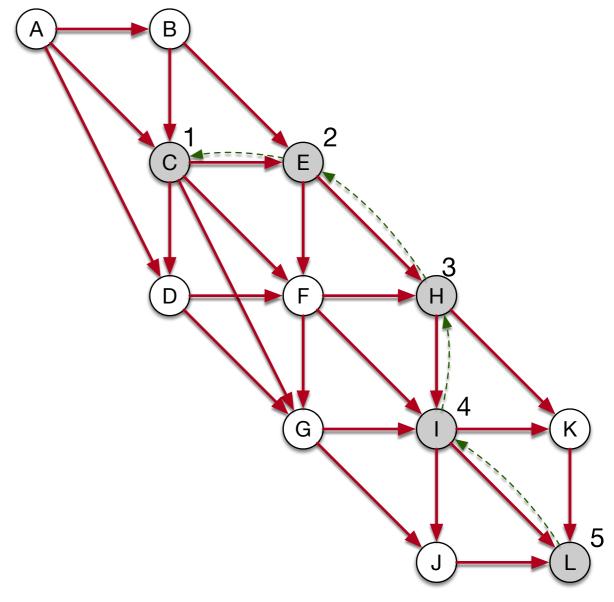
dfs_visit(L)

dfs_visit(I)

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



we pick v = L and call
dfs\_visit(L)
this colors L gray

```
OS stack

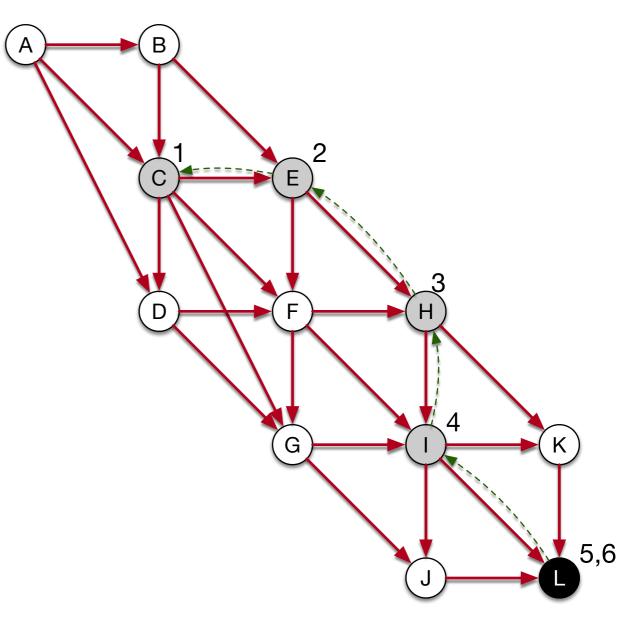
dfs_visit(L)

dfs_visit(I)

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



L has no vertices in the adjacency list Therefore, we finally go to the last line

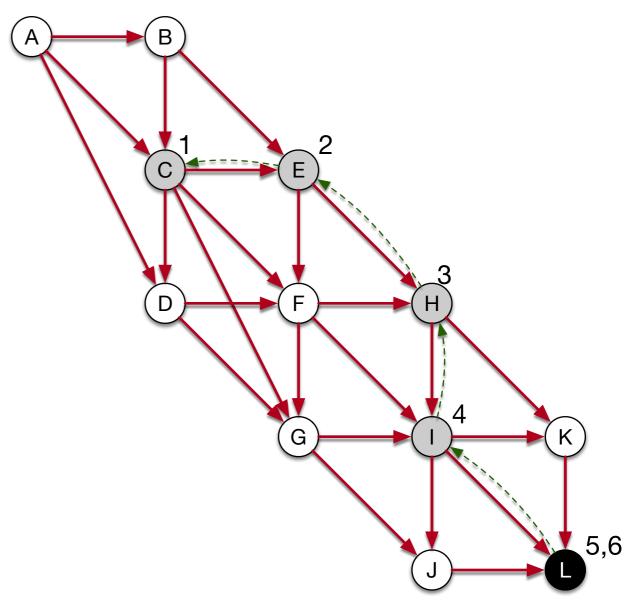
```
OS stack

dfs_visit(I)

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



We finish dfs\_visit(L) and are back at the execution of dfs\_visit(I)

```
OS stack

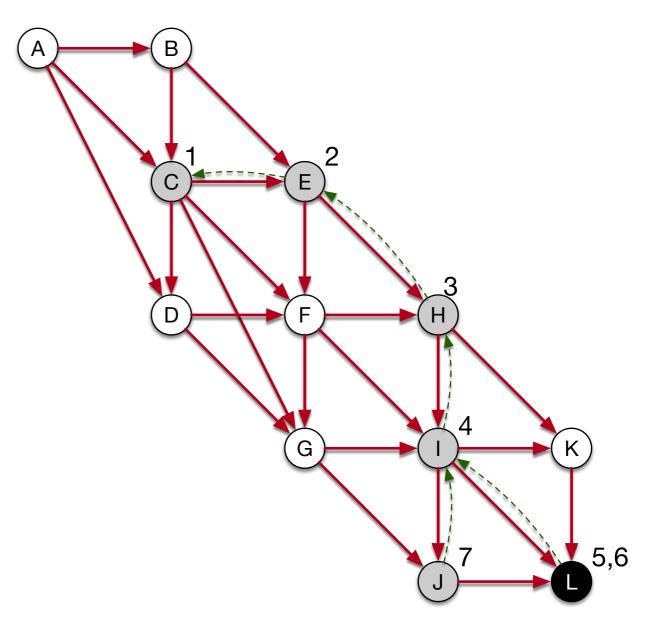
dfs_visit(J)

dfs_visit(I)

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



We pick a white vertex reachable from I: J

```
OS stack

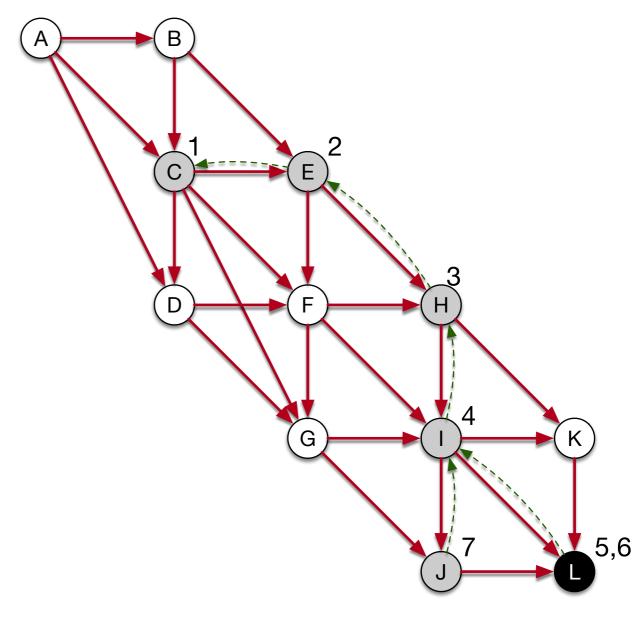
dfs_visit(J)

dfs_visit(I)

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



There are no white nodes in the adjacency list of J

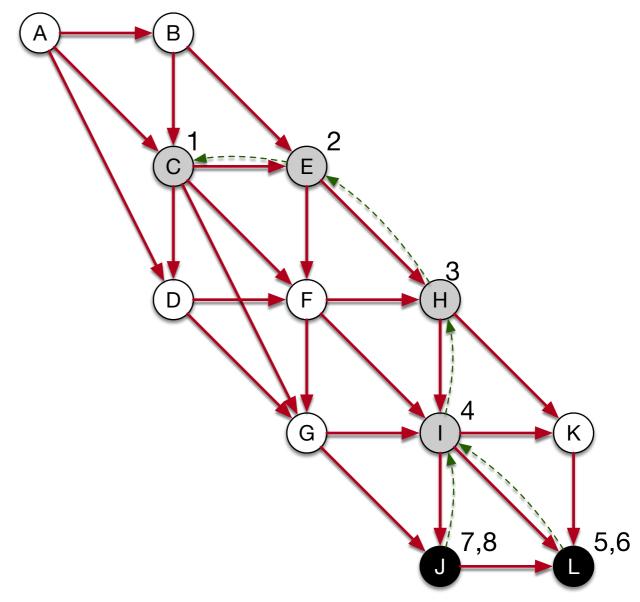
```
OS stack

dfs_visit(I)

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



We close the call on J and are back to dfs\_visit(I)

```
OS stack

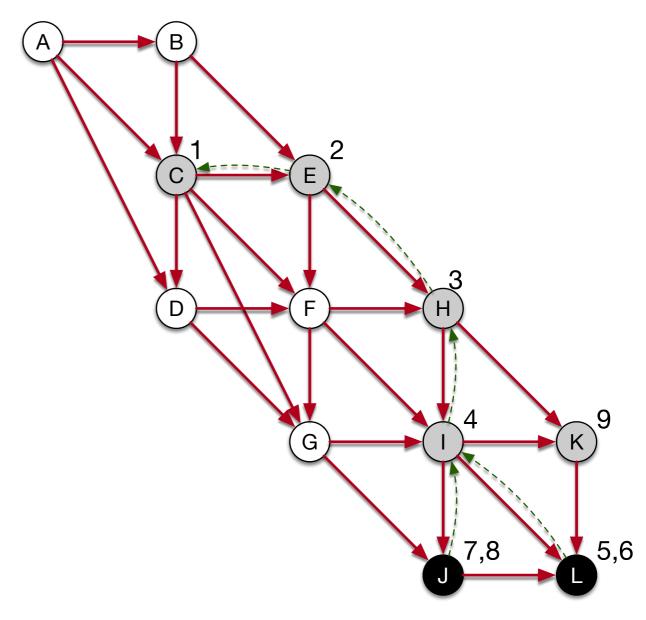
dfs_visit(K)

dfs_visit(I)

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



dfs\_visit(I) now goes to K

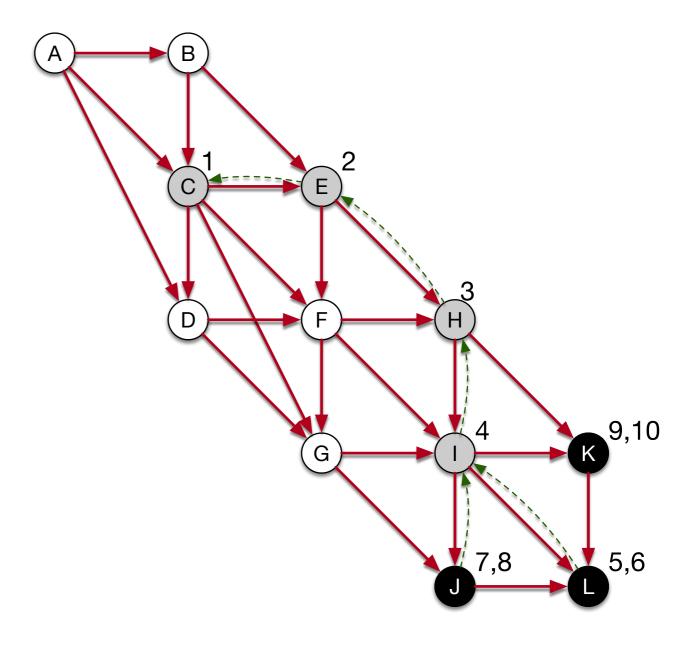
```
OS stack

dfs_visit(I)

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



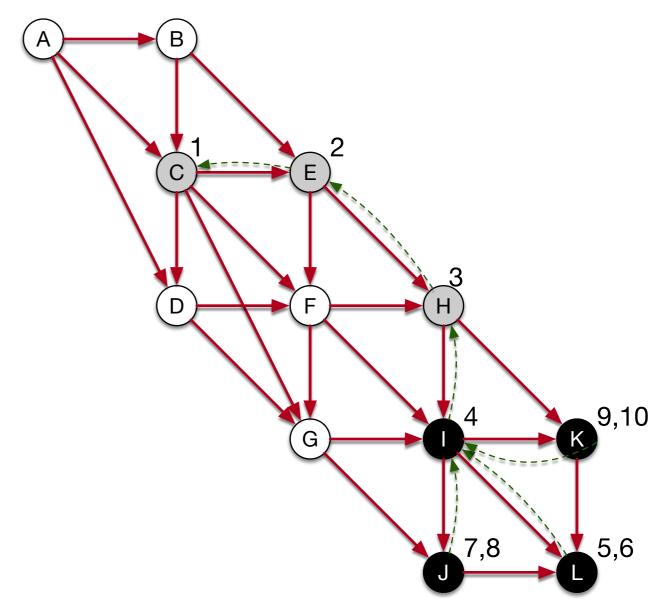
dfs\_visit(K) finishes

```
OS stack

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```



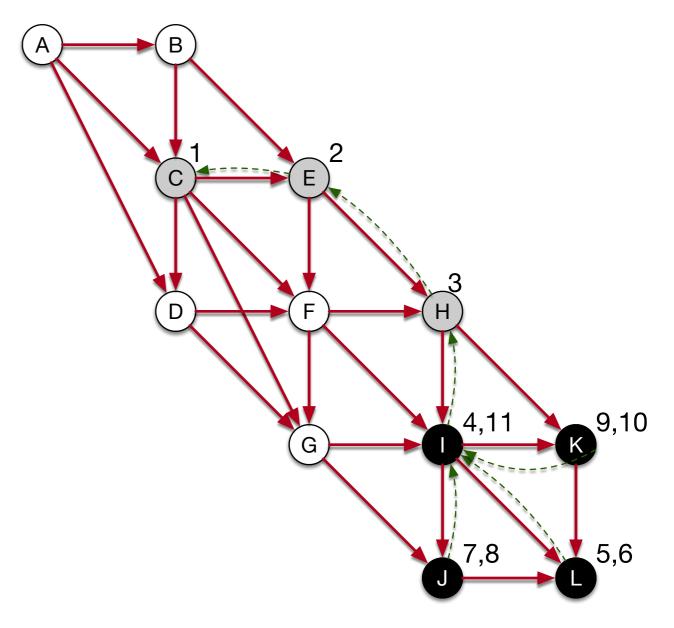
dfs\_visit(I) runs again
but finds no white vertices,
so it finishes

```
OS stack

dfs_visit(H)

dfs_visit(E)

dfs_visit(C)
```

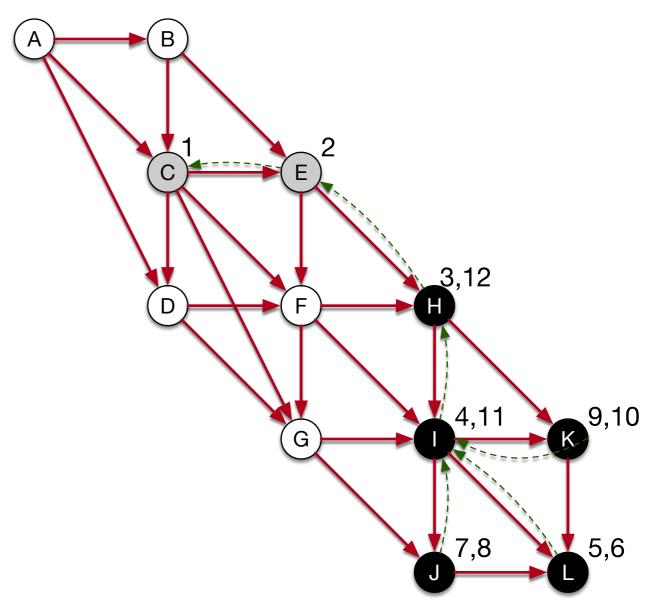


dfs\_visit(I) runs again
but finds no white vertices,
so it finishes

```
OS stack

dfs_visit(E)

dfs_visit(C)
```

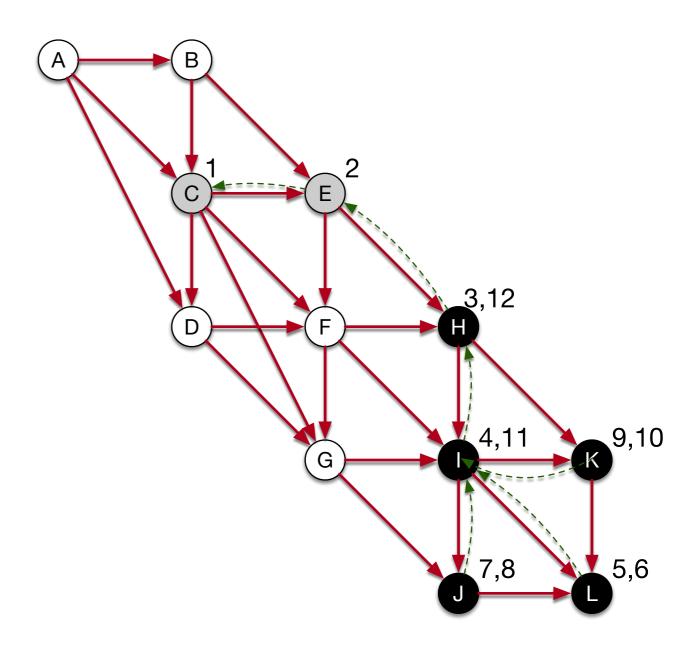


dfs\_visit(H) runs again
but finds no white vertices,
so it finishes

```
OS stack

dfs_visit(E)

dfs_visit(C)
```



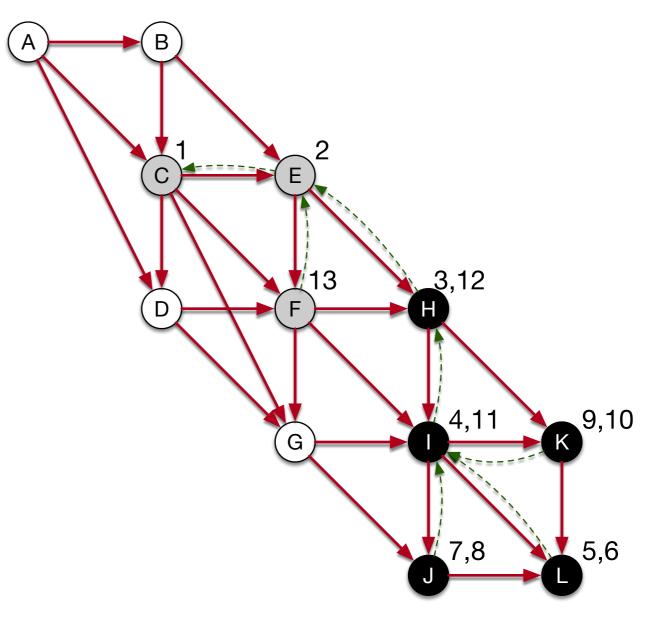
dfs\_visit(E) runs again

```
OS stack

dfs_visit(F)

dfs_visit(E)

dfs_visit(C)
```



Finds F

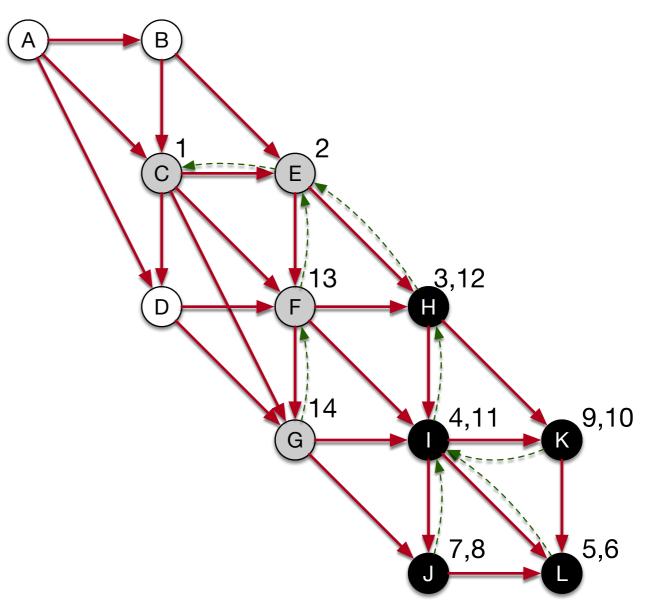
```
OS stack

dfs_visit(G)

dfs_visit(F)

dfs_visit(E)

dfs_visit(C)
```



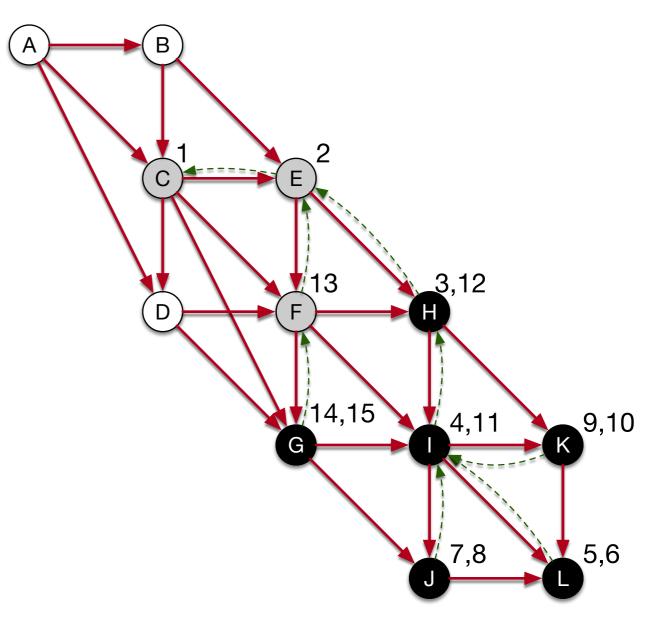
Finds G

```
OS stack

dfs_visit(F)

dfs_visit(E)

dfs_visit(C)
```

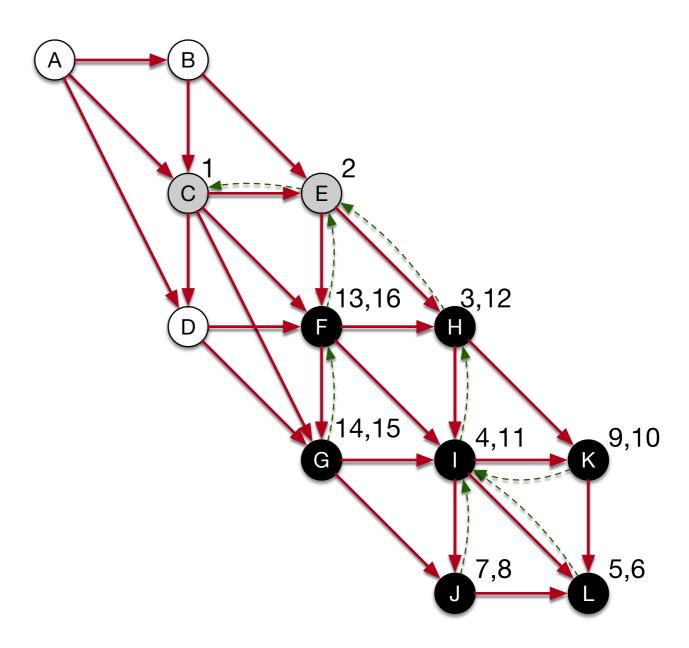


Nothing left in G

```
OS stack

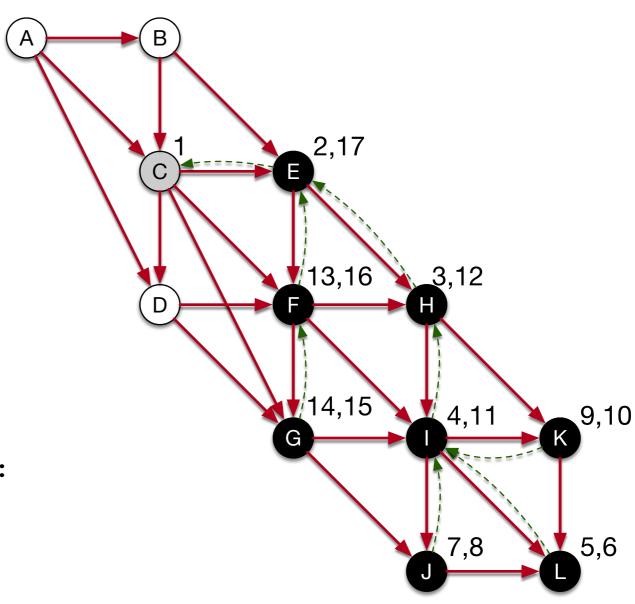
dfs_visit(E)

dfs_visit(C)
```



Nothing left in F

```
OS stack
dfs_visit(C)
```

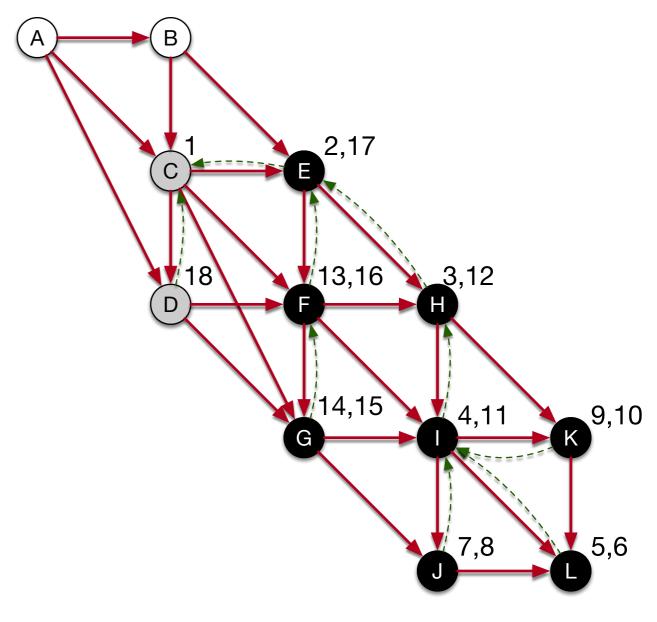


Finishing E

```
OS stack

dfs_visit(D)

dfs_visit(C)
```

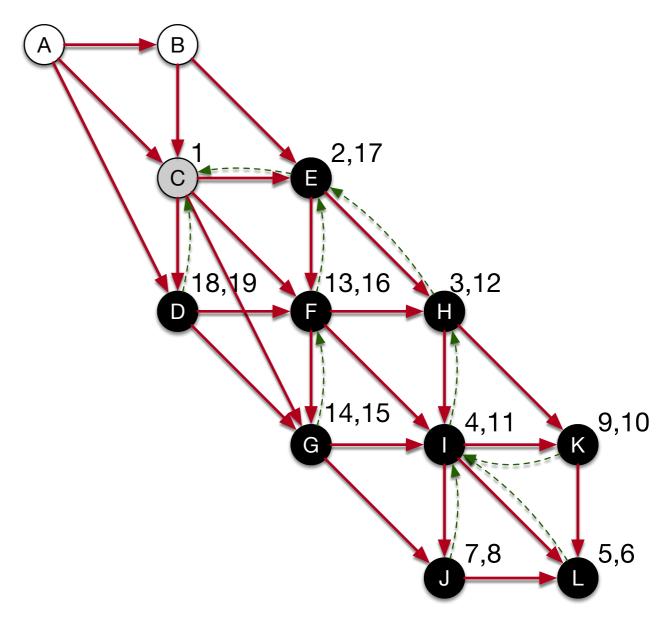


Last white vector in the adjacency list of C is D

```
OS stack

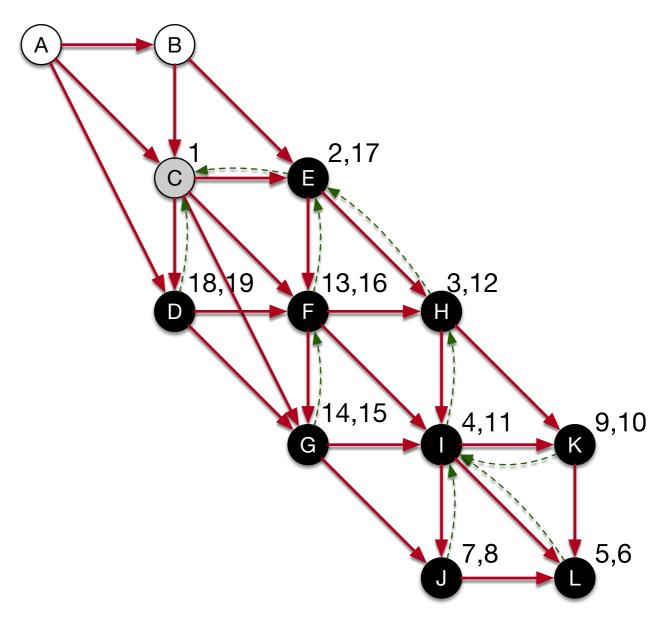
dfs_visit(D)

dfs_visit(C)
```



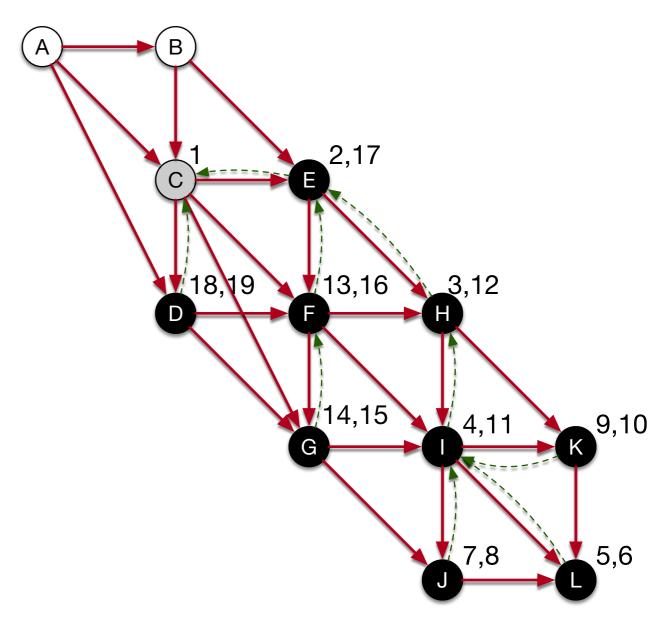
D has no white vertices in its adjacency list

```
OS stack dfs_visit(C)
```



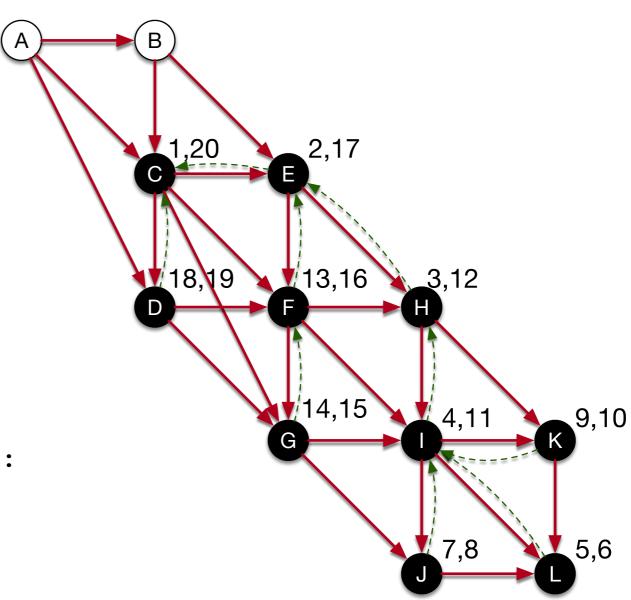
We are back in C

```
OS stack dfs_visit(C)
```



Now we close C

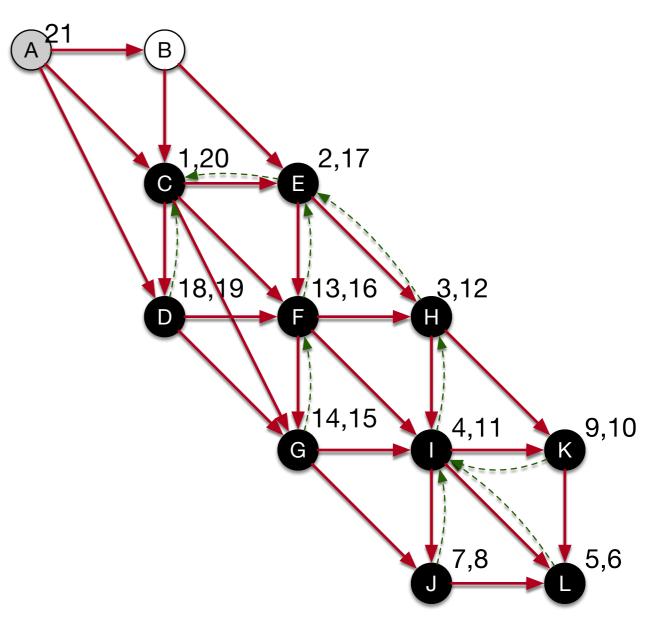
```
OS stack
```



Now we close C

- At this point, the original call to dfs\_visit(C) is done
- However, since there are still white nodes left, we have to pick one of them and visit again.
- We pick A

```
OS stack dfs_visit(A)
```

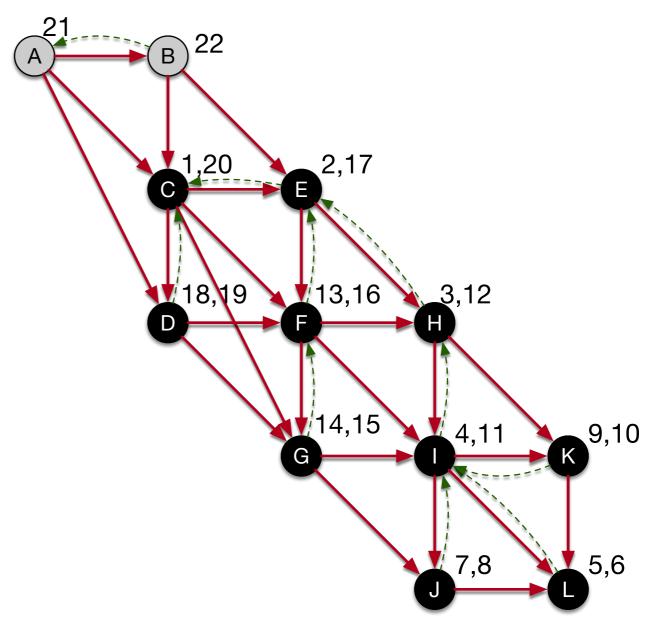


A is the only node in the stack

```
OS stack

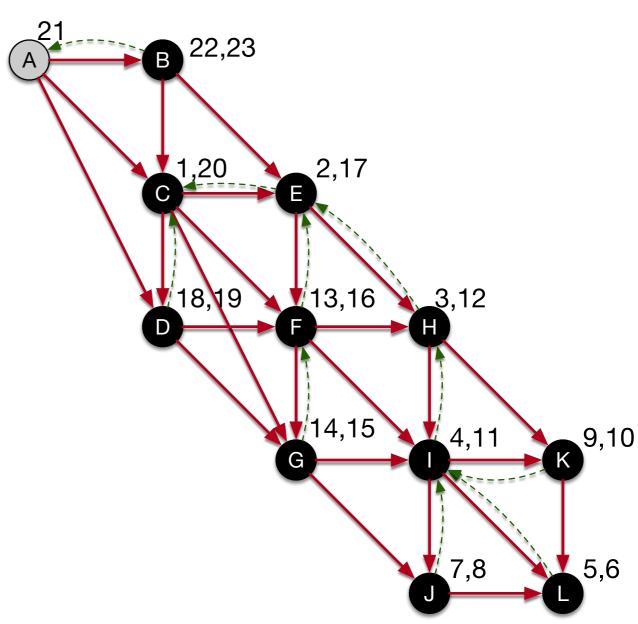
dfs_visit(B)

dfs_visit(A)
```



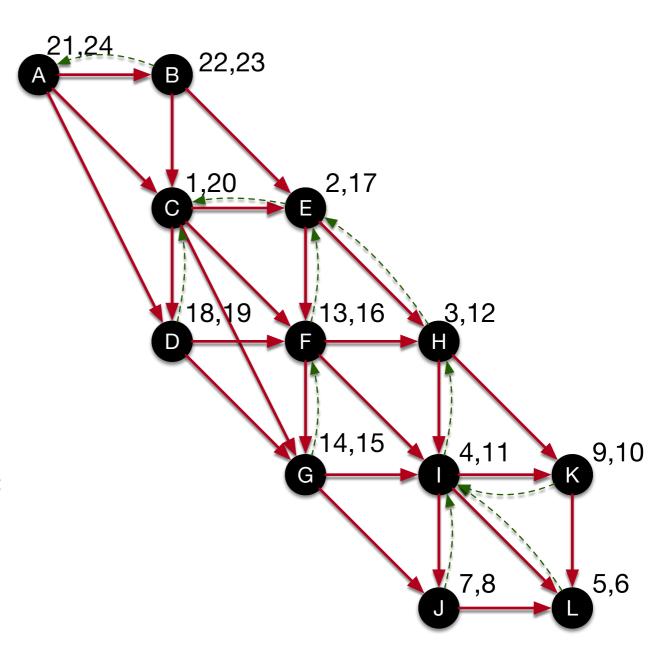
We discover B from A

```
OS stack dfs_visit(A)
```



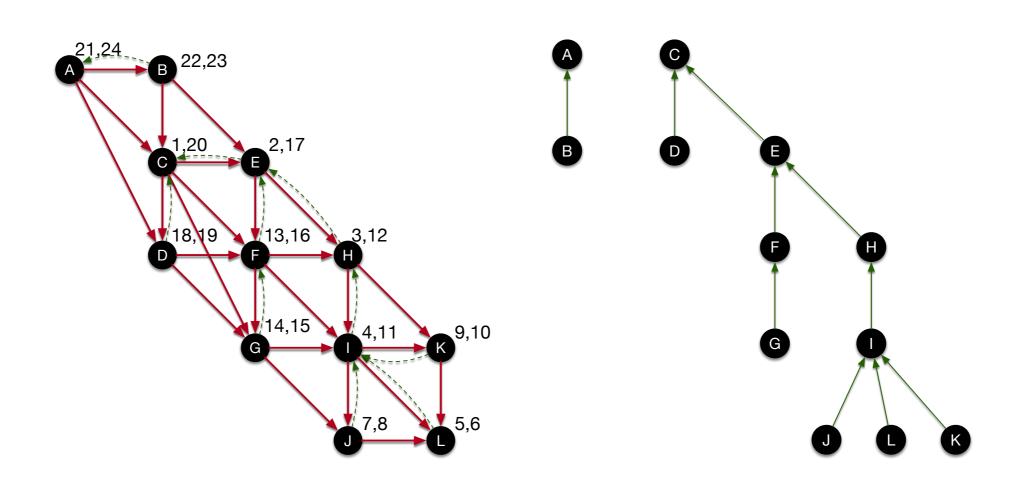
We can finish B

```
OS stack dfs_visit(A)
```



We can finish A

- Now we are done
  - The predecessor relationship has given us a nice set of trees — a "forest"

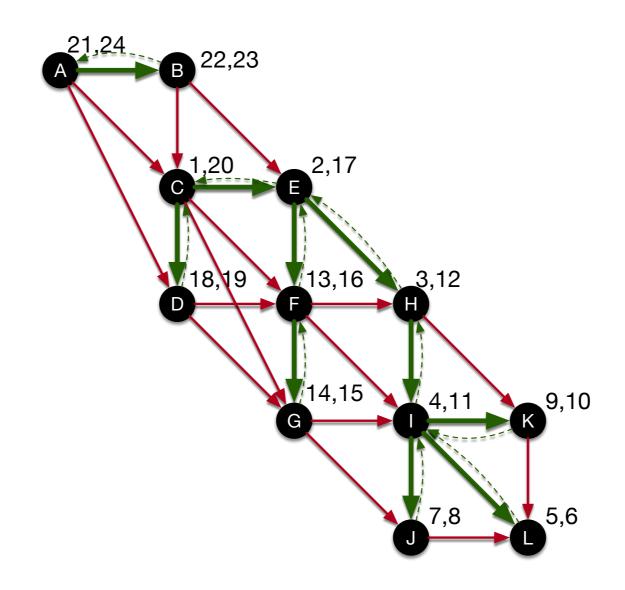


- Runtime of algorithm
  - We look at all the elements of the adjacency lists
  - For each, we do constant work
  - But we also need to do some initial work for all vertices
  - Runtime is  $\Theta(\max(|V|, |E|))$

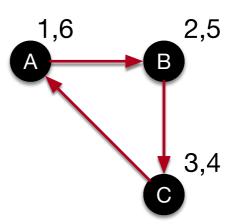
- Properties:
  - Parenthesis Theorem
    - If for two nodes
      - If  $[u.d,u.f] \cap [v.d,v.f] = \emptyset$  then neither u and v are descendants in the predecessor forest
      - If  $[u.d,u.f] \subset [v.d,v.f]$  then u is a descendant of v
      - If  $[u.d,u.f] \supset [v.d,v.f]$  then v is a descendant of u

- White Path Theorem
  - v is a descendant of u exactly if
  - At the time of discovery of u there is a path from u to v consisting entirely of white vertices

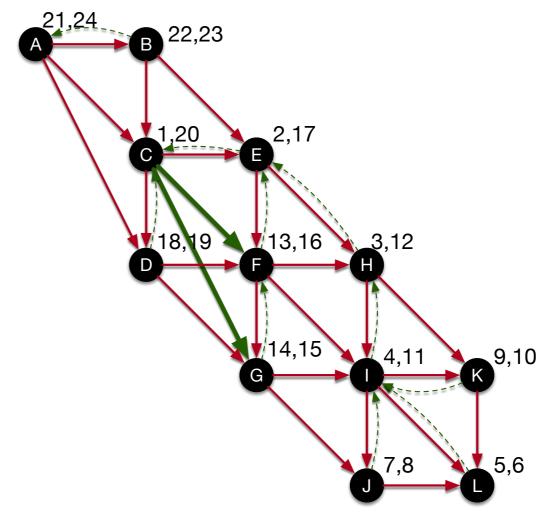
- Classification of edges:
  - Tree edges are edges in the depth first tree



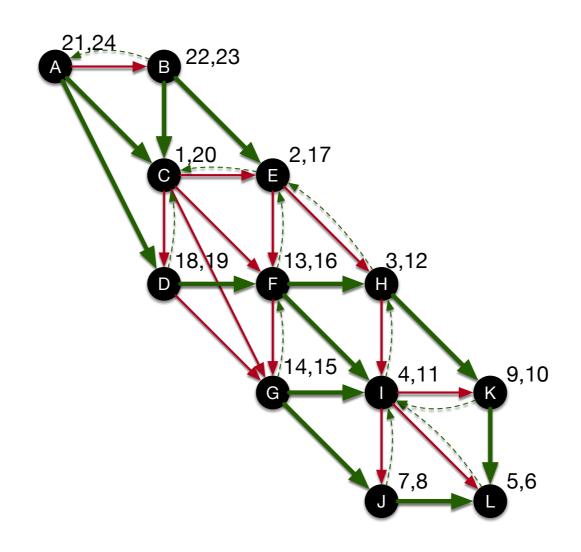
- Back edges are edges go from a descendant to an ancestor
  - Simple example:
    - Start in A, discover B, discover C
    - Edge from C to A is a back edge



- Forward edges
  - Edges connecting an ancestor to a descendant, but that are not in the tree



- Cross Edges: anything else
  - Can be in the same tree or connecting different trees



- If we look at an edge (u,v) during depth first search for the first time
  - (In an undirected graph, we look at each edge twice)
  - If v is white: tree edge
  - If v is gray: back edge
  - If v is black: forward or cross edge

- In a depth first search on an <u>undirected</u> graph, every edge is either a tree edge or a back edge
  - Let (u, v) be an edge and assume that u is discovered first:  $u \cdot d < v \cdot d$
  - The algorithm discovers and finishes v before u, so  $u \cdot f > v \cdot f$
  - If DFS uses the edge (u, v) from u, then v is white, and (u, v) becomes a tree edge
  - If DFS uses the edge (u, v) from v, then u is gray at this moment and this becomes a back edge.