## Activity: Linear Hashing

## File State

- The file state of an LH file is determined by the number $n$ of buckets
- level $l$
- split pointer $S$
- Formula: $n=2^{l}+s$ with $l$ as high as possible, i.e. with $s \in\left\{0,1, \ldots, 2^{l}-1\right\}$


## File State

- Clarification regarding the literature
- The original LH scheme can start with any number of buckets
- In this class, we are using the most common case


## Exercise

- What is the level and the state of an LH file with 13 buckets?


## Solution

- We write $13=2^{3}+5$
- Level is $l=\left\lfloor\log _{2}(13)\right\rfloor=3$
- Split pointer is $s=13-l$


## Exercise

- Where would the records with the following (randomly picked) keys be inserted?
- 82
- 27
- 37


## Solution

- Level is 3 , so we use first remainder modulo $2^{3}=8$ and $2^{4}=16$ second
- $82(\bmod 8)=2$. Since $2<5$, we rehash: $82(\bmod 16)=2$ and we insert into bucket 2
- $27(\bmod 8)=3$. Since $3<5$, we rehash:
$27(\bmod 16)=11$. We insert into bucket 11
- $37(\bmod 8)=5$. Since $5 \nless 5$, we do not rehash but insert into bucket 5 .


## Exercise

- Where would the records with the following (randomly picked) keys be inserted?
- 48
- 60
- 63
- 71


## Solution

- $48(\bmod 8)=0$. Rehash: $48(\bmod 16)=0$ and insert into bucket 0 .
- $60(\bmod 8)=4$. Rehash: $60(\bmod 16)=12$ and insert into bucket 12.
- $63(\bmod 8)=7$. Rehash not necessary. Insert into bucket 7.
- $71(\bmod 8)=7$. No rehash is necessary.


## Exercise

- Where would the records with the following (randomly picked) keys be inserted?
- 98
- 75
- 25
- 30


## Solution

- $98(\bmod 8)=2$. Rehash: $98(\bmod 16)=2$. Insert into bucket 2
- $75(\bmod 8)=3$. Rehash: $75(\bmod 16)=11$. Insert into bucket 11
- $25(\bmod 8)=1$. Rehash: $25(\bmod 16)=9$. Insert into bucket 9 .
- $30(\bmod 8)=6$. Insert into bucket 6 .


## Exercise

- Give the level and split pointer values as an LH file moves from 6 buckets to 20


## Solution

| Nr o Buckets | Level | Split Ptr |
| :---: | :---: | :---: |
| 6 | 2 | 2 |
| 7 | 2 | 3 |
| 8 | 3 | 0 |
| 9 | 3 | 1 |
| 10 | 3 | 2 |
| 11 | 3 | 3 |
| 12 | 3 | 4 |
| 13 | 3 | 5 |
| 14 | 3 | 6 |
| 15 | 3 | 7 |
| 16 | 4 | 0 |
| 17 | 4 | 1 |
| 18 | 4 | 2 |
| 19 | 4 | 3 |
| 20 |  | 4 |

## Interpretation

- We can encapsulate the behavior of the level and split pointer into the following algorithm

```
def split(level, split_pointer):
    split pointer += 1
    if split_pointer == 2**level:
        split_pointer = 0
        level += 1
    return (level, split_pointer)
```

- We increment the split pointer
- If the split pointer equals $2^{\text {level }}$ then set the split pointer to zero and increment the level


## Programming Exercise

- Using a programming platform of your choice, implement the LH addressing algorithm
- Insert 1000 records with key uniformly selected between 0 and $2^{32}-1$ into an LH file with (a) 12 and (b) 25 buckets.
- Look at the size of the buckets.


## Solution

- I changed the number to $1,000,000$
- For 12 buckets:



## Solution

- Here is the chart for 25 buckets



## Interpretation

- Even with a perfect hash function, an LH file has buckets of equal size only if the number of buckets is a power of two.
- Otherwise, there are buckets already split in the current round and those not yet split.
- The latter have about twice as many records

